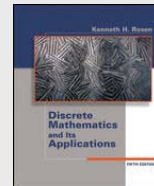


Chapter 10 (Part 1): Boolean Algebra

- ✓ Boolean Functions (10.1)
- ✓ Representing Boolean Functions (10.2)
- ✓ Logic Gates (10.3)



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- ✓ It has started from the book titled “The laws of thought” written by George Boole in 1854
- ✓ Claude Shannon showed how the basic rules of logic could be used to design circuits
- ✓ The circuit in electronic devices have inputs, each of which is either 0 or 1 and produce outputs that are also 0 or 1

- ✓ The operation of a circuit is defined by a Boolean function that specifies the value of an output for each set of inputs
- ✓ One of the goal is to describe methods for finding a simplified expression (min number of sums and products) that represents a Boolean function (Karnaugh maps).

Boolean Functions (10.1)

- ✓ Introduction
 - Electronic and optimal switches are studied using the set $\{0, 1\}$ and the rules of Boolean algebra
 - The 3 operations in Boolean algebra that are used are Complementation, Boolean Sum and the Boolean Product

Boolean Functions (10.1) (cont.)

- Complement of an element is defined by:

$$\bar{0} = 1; \bar{1} = 0$$

- Boolean sum denoted by + or by OR obeys:

$$1 + 1 = 1, \quad 1 + 0 = 1, \quad 0 + 1 = 1, \quad 0 + 0 = 0$$

- Boolean product denoted by AND obeys:

$$1 \cdot 1 = 1, \quad 1 \cdot 0 = 0, \quad 0 \cdot 1 = 0, \quad 0 \cdot 0 = 0.$$

Boolean Functions (10.1) (cont.)

- Example: Find the value $1 \cdot 0 + \overline{(0 + 1)}$

$$\text{Solution: } 1 \cdot 0 + \overline{(0 + 1)} = 0 + \bar{1} = 0 + 0 = 0$$

✓ Boolean expressions & Boolean Functions

- Definition

Let $B = \{0, 1\}$ then $B^n = \{(x_1, x_2, \dots, x_n) \text{ such that } x_i \in B \text{ for } 1 \leq i \leq n\}$. The variable x is called a Boolean variable if it assumes values only from B (0 or 1). A function B^n to B is called a Boolean function of degree n .

Boolean Functions (10.1) (cont.)

- Example: The function $F(x, y) = x\bar{y}$ from $B^2 = \{(x, y) \text{ such that } (x, y) \in \{0, 1\} * \{0, 1\}\}$ is a Boolean function with:

$$F(1, 1) = 0, F(1, 0) = 1, F(0, 1) = 0, F(0, 0) = 0$$

- Example: Find the values of the Boolean function represented by $F(x, y, z) = xy + \bar{z}$

Solution: The values of this function are displayed on the following table:

Boolean Functions (10.1) (cont.)

x	y	z	xy	\bar{z}	$F(x, y, z) = xy + \bar{z}$
1	1	1	1	0	1
1	1	0	1	1	1
1	0	1	0	0	0
1	0	0	0	1	1
0	1	1	0	0	0
0	1	0	0	1	1
0	0	1	0	0	0
0	0	0	0	1	1

✓ Identities of Boolean algebra

Identity	Name
$\overline{\overline{x}} = x$	Law of the double component
$x + x = x$ $x \cdot x = x$	Idempotent laws
$x + 0 = x$ $x \cdot 1 = x$	Identity laws
$x + 1 = 1$ $x \cdot 0 = 0$	Domination laws
$x + y = y + x$ $xy = yx$	Commutative laws
$x + (y + z) = (x + y) + z$ $x(yz) = (xy)z$	Associative laws
$x + yz = (x + y)(x + z)$	Distributive laws
$\overline{(xy)} = \overline{x} + \overline{y}$ $\overline{(x + y)} = \overline{x} \overline{y}$	De Morgan's laws
$x + xy = x$ $x(x + y) = x$	Absorption laws
$x + \overline{x} = 1$	Unit property
$x\overline{x} = 0$	Zero property

Boolean Functions (10.1) (cont.)

– Example: Show that the distributive law
 $x(y + z) = xy + xz$

Solution:

x	y	z	y + z	xy	xz	x(y + z)	xy + xz
1	1	1	1	1	1	1	1
1	1	0	1	1	0	1	1
1	0	1	1	0	1	1	1
1	0	0	0	0	0	0	0
0	1	1	1	0	0	0	0
0	1	0	1	0	0	0	0
0	0	1	1	0	0	0	0
0	0	0	0	0	0	0	0

Boolean Functions (10.1) (cont.)

✓ Duality

– Goal: the dual of a Boolean expression is obtained by interchanging Boolean sums and Boolean products and interchanging 0s and 1s

– Example: Find the duals of $x(y + 0)$ and

$$\bar{x}.1 + (\bar{y} + z)$$

Solution: the duals are $x + (y.1)$ and $(\bar{x} + 0)(\bar{y}z)$

Boolean Functions (10.1) (cont.)

– The duality principle enables to obtain more identities. Since an identity between functions represented by Boolean expressions remains valid when the duals of both sides of the identity are taken

✓ The abstract definition of a Boolean Algebra

– Goal: Recognize that a particular structure (set, propositions,...) is a Boolean algebra \Rightarrow all results of Boolean algebra will apply

Boolean Functions (10.1) (cont.)

– Definition 1:

A Boolean algebra is a set B with two binary operations \vee and \wedge , elements 0 and 1 , and a unary operation $\bar{}$ such that these properties hold for all x, y and z in B :

$$\left. \begin{array}{l} x \vee 0 = x \\ x \wedge 1 = x \end{array} \right\} \textit{Identity laws}$$

$$\left. \begin{array}{l} x \vee \bar{x} = 1 \\ x \wedge \bar{x} = 0 \end{array} \right\} \textit{Complement laws}$$

$$\left. \begin{array}{l} (x \vee y) \vee z = x \vee (y \vee z) \\ (x \wedge y) \wedge z = x \wedge (y \wedge z) \end{array} \right\} \textit{Associative laws}$$

$$\left. \begin{array}{l} x \vee y = y \vee x \\ x \wedge y = y \wedge x \end{array} \right\} \textit{Commutative laws}$$

$$\left. \begin{array}{l} x \vee (y \wedge z) = (x \vee y) \wedge (x \vee z) \\ x \wedge (y \vee z) = (x \wedge y) \vee (x \wedge z) \end{array} \right\} \textit{Distributive laws}$$

Boolean Functions (10.1) (cont.)

- The set of propositions in n variables, with the \vee and \wedge operators, \mathbf{F} and \mathbf{T} , and the negation operator, also satisfies all the properties of a Boolean algebra
- Similarly, the set of subsets of a universal set Ω with the union and intersection operators, the empty set and the universal set, and the set complementation operator is a Boolean algebra

Representing Boolean Functions (10.2)

- ▼ Two problems of Boolean algebra are emphasized in this section
 1. Given the values of a Boolean function, how can a Boolean expression that represents this function be found?
 2. Is there a smaller set of operators that can be used to represent all Boolean functions?

Representing Boolean Functions (10.2) (cont.)

- ▼ Sum-of-products expansions
 - Goal: Find a Boolean expression that represents a Boolean function
 - Example: Find Boolean expressions that represent the functions $F(x, y, z)$ and $G(x, y, z)$ given by the following table:

x	y	z	F	G
1	1	1	0	0
1	1	0	0	1
1	0	1	1	0
1	0	0	0	0
0	1	1	0	0
0	1	0	0	1
0	0	1	0	0
0	0	0	0	0

Representing Boolean Functions (10.2) (cont.)

x	y	z	F	G
1	1	1	0	0
1	1	0	0	1
1	0	1	1	0
1	0	0	0	0
0	1	1	0	0
0	1	0	0	1
0	0	1	0	0
0	0	0	0	0

Solution:

- To form the Boolean expressions for the function $F(x,y,z)$, we need to look at the table and notice that F has the value 1 only if $x = 1$ and $y = 0$ and $z = 1$ otherwise F has value 0
 $\Rightarrow F(x,y,z) = x\bar{y}z$
- To form the Boolean expression for the function $G(x,y,z)$, we have to notice that $G(x,y,z) = 1 \Leftrightarrow x = y = 1$ and $z = 0$ or $x = z = 0$ and $y = 1$; otherwise $G(x,y,z) = 0$
 $\Rightarrow G(x,y,z) = (xy\bar{z}) + (\bar{x}y\bar{z})$

Representing Boolean Functions (10.2) (cont.)

– Definition 1:

A literal is a Boolean variable or its complement. A minterm of the Boolean variables x_1, x_2, \dots, x_n is a Boolean product $y_1y_2\dots y_n$, where $y_i = x_i$ or $y_i = \bar{x}_i$. Hence, a minterm is a product of n literals, with one literal of each variable.

The minterm $y_1y_2\dots y_n$ is 1 if and only if each y_i is 1, and therefore $x_i = 1$ when $y_i = x_i$ and $x_i = 0$ when $y_i = \bar{x}_i$.

Representing Boolean Functions (10.2) (cont.)

- Example: Find a minterm that equals 1 if $x_1 = x_3 = 0$ and $x_2 = x_4 = x_5 = 1$, and equals 0 otherwise

Solution: The minterm is: $\bar{x}_1 x_2 \bar{x}_3 x_4 x_5$

- Minterms are used to determine the Boolean expression of a Boolean function in a table

Representing Boolean Functions (10.2) (cont.)

- Example: Find the sum-of-products expansion for the function $F(x, y, z) = (x + y)\bar{z}$

Solution (1):

$$\begin{aligned}(x + y)\bar{z} &= x\bar{z} + y\bar{z} = x1\bar{z} + 1y\bar{z} \\ &= x(y + \bar{y})\bar{z} + (x + \bar{x})y\bar{z} \\ &= xy\bar{z} + x\bar{y}\bar{z} + xy\bar{z} + \bar{x}y\bar{z} \\ &= xy\bar{z} + x\bar{y}\bar{z} + \bar{x}y\bar{z} \quad (\text{since } u + u = u)\end{aligned}$$

Representing Boolean Functions (10.2) (cont.)

Solution (2): We use the minterm technique from a table

x	y	z	$x + y$	\bar{z}	$(x + y) \bar{z}$
1	1	1	1	0	0
1	1	0	1	1	1
1	0	1	1	0	0
1	0	0	1	1	1
0	1	1	1	0	0
0	1	0	1	1	1
0	0	1	0	0	0
0	0	0	0	1	0

Representing Boolean Functions (10.2) (cont.)

We need only to spot the set of values (x,y,z) for which

This set is: $F(x, y, z) = (x + y) \bar{z} = 1$.

$\{(1,1,0);(1,0,0);(0,1,0)\}$

We need to find the minterm that equals 1 if

$$\underbrace{(x = y = 1 \text{ and } z = 0)}_{xy\bar{z}} \text{ OR } \underbrace{(x = 1 \text{ and } y = z = 0)}_{x\bar{y}\bar{z}} \text{ OR } \underbrace{(x = z = 0 \text{ and } y = 1)}_{\bar{x}y\bar{z}}$$

- Theorem 1: Every Boolean function can be represented using the three Boolean operators \cdot , $+$, and $\bar{}$ (complement)

Representing Boolean Functions (10.2) (cont.)

✓ Functional completeness

- The set $\{., +, '\}$ is functionally complete because of theorem 1
- Question: Can we find a smaller set of functionally complete operators?
- We answer “yes” to this question if one of the three operators can be expressed in terms of the other two
- One of the De Morgan’s laws provide the solution:
$$x + y = \overline{\overline{x} \overline{y}} \text{ since } \overline{(x + y)} = \overline{x} \overline{y}$$

Representing Boolean Functions (10.2) (cont.)

- Similarly, we have: $xy = \overline{\overline{x} + \overline{y}}$
- Finally, we have found a smaller set functionally complete; it contains only 2 operators!
- Question: Can we still reduce the set to only one operator? The answer is “yes”.
- Indeed, there are 2 sets that contain only one operator, they are: the **NAND** and the **NOR**.

Representing Boolean Functions (10.2) (cont.)

- NAND ($|$) is defined as

$$1|1 = 0 \text{ and } 1|0 = 0|1 = 0|0 = 1$$

- NOR (\downarrow) is defined as

$$1\downarrow 1 = 1\downarrow 0 = 0\downarrow 1 = 0 \text{ and } 0\downarrow 0 = 1$$

- Since the set $\{., '\}$ is functionally complete and

$\bar{x} = x|x$, $xy = (x|y) | (x|y)$ then the set $\{| \}$ is functionally complete.

- Home-exercise: Prove that the set $\{\downarrow\}$ is functionally complete.

Logic Gates (10.3)

▼ Introduction

- A computer, or electronic devices, is made up of a number of circuits
- Each circuit can be designed using the rules of Boolean algebra
- The basic element of circuits are called gates
- Each type of gates implements a Boolean Operation

Logic Gates (10.3) (cont.)

- All circuits studied in this chapter provides output that depends only on the input, and not on the current state of the circuit
- These circuits that have no memory capabilities are called combinational circuits or gating networks
- Combinatorial circuits are built using 3 types of elements: (1) inverter, (2) OR gate, and (3) AND gate

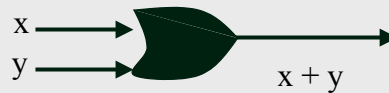
Logic Gates (10.3) (cont.)

- Role of each type of elements:
 1. Inverter: it accepts the value of one Boolean variable as input and produces the complement of this value as its output.



Logic Gates (10.3) (cont.)

- OR gate: it accepts as input values of 2 or more Boolean variables. The output is the Boolean sum of their values



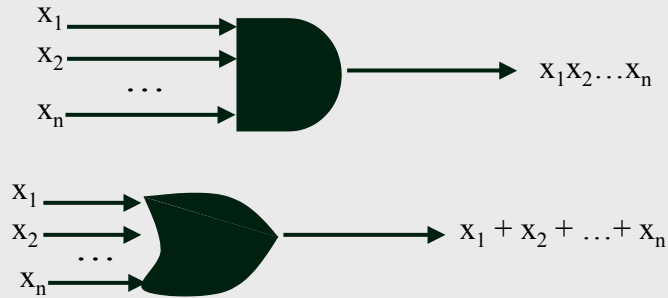
Logic Gates (10.3) (cont.)

- AND gate: The inputs to this gate are the values of two or more Boolean variables. The output is the Boolean product of their values.



Logic Gates (10.3) (cont.)

- We can have multiple inputs to AND and OR:

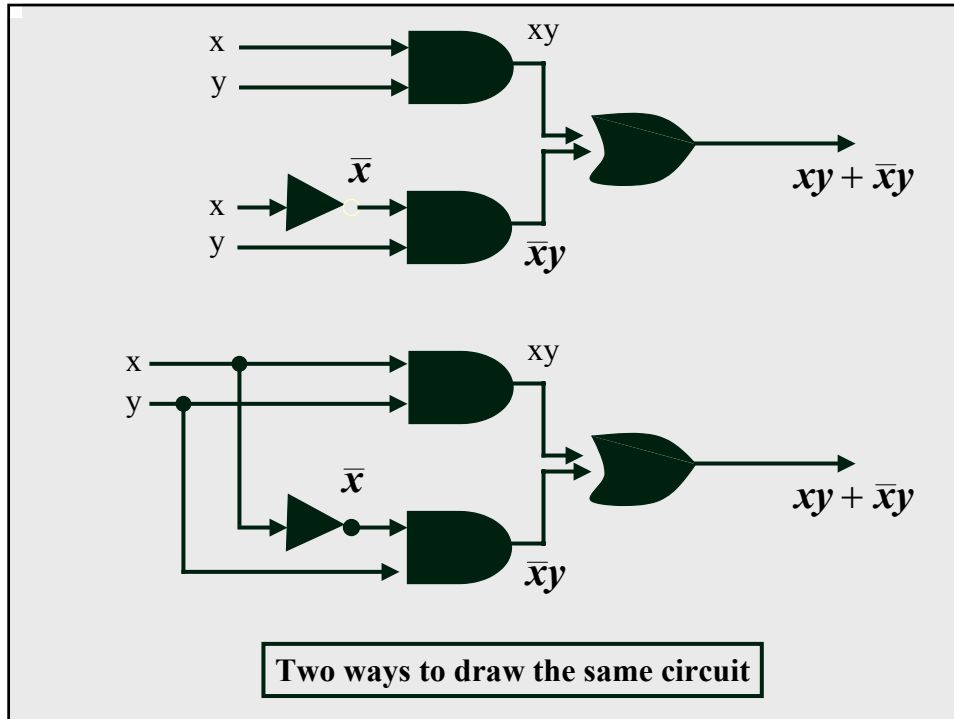


Gates with n inputs

Logic Gates (10.3) (cont.)

✓ Combinations of Gates

- Combinational circuits can be constructed using a combination of inverters, OR gates, and AND gates
- When combinations of circuits are formed, some gates may share input using branching
- The following figure shows 2 ways of drawing the same circuit



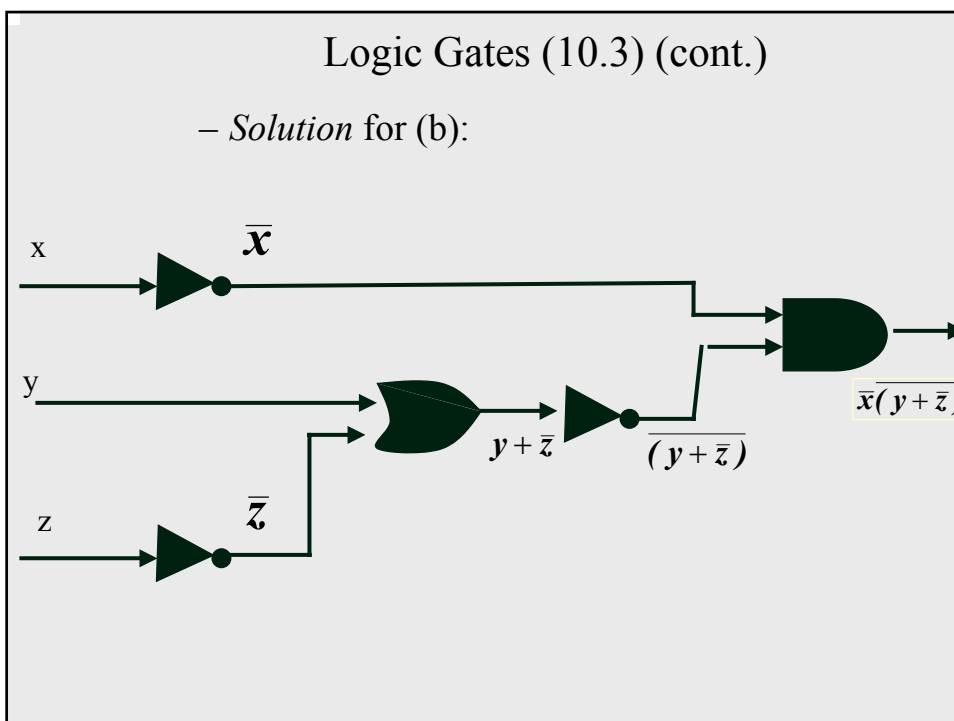
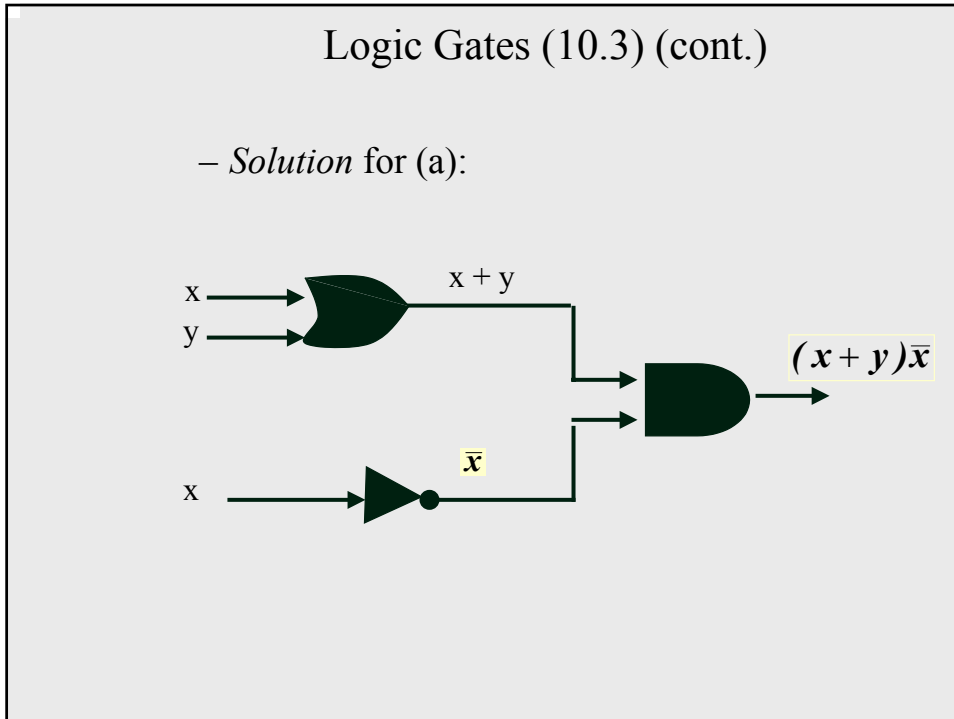
Logic Gates (10.3) (cont.)

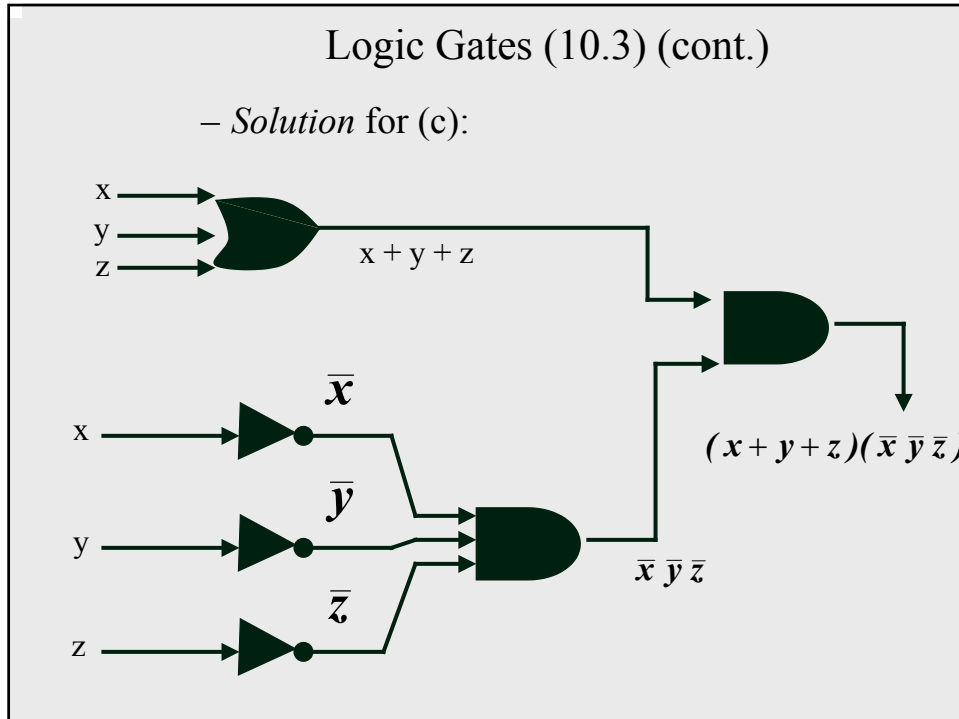
– Example: Construct circuits that produce the following outputs:

a) $(x + y)\bar{x}$

b) $\bar{x}(y + \bar{z})$

c) $(x + y + z)(\bar{x}\bar{y}\bar{z})$





Logic Gates (10.3) (cont.)

✓ Example of circuits

- Example: A committee of 3 individuals decides issues for an organization. Each individual votes either yes or no for each proposal that arises. A proposal is passed if it receives at least 2 yes votes. Design a circuit that determines whether a proposal passes.

Solution: We want to have the following Boolean function: $xy + xz + yz$ represented by the following circuit:

